

# An algebraic algorithm for blind source separation and tensor decomposition

Martina Iannacito

joint work with Ignat Domanov and Lieven De Lathauwer

New directions in linear algebra  
BIRS, Canada, August 31, 2023



# Overview

## The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

## The Canonical Polyadic Decomposition

From the theorem to the algorithm

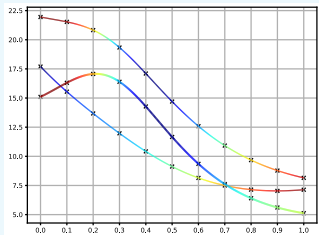
The bottleneck

## Algorithm improvements

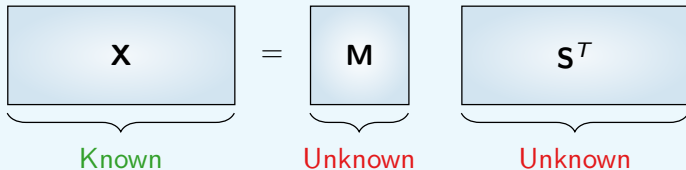
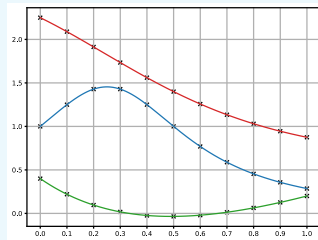
Exterior algebra

Numerical results

# Blind Source Separation problem



$$= \begin{bmatrix} 8 & 1 & 6 \\ 3 & 5 & 7 \\ 4 & 9 & 2 \end{bmatrix} \times$$



## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

Algorithm improvements

Exterior algebra

Numerical results

## Constraints for uniqueness

**Definition:** A *deterministic* condition on  $\mathbf{X}$  imposes a particular property of  $\mathbf{X}$  that is always true.

---

**Definition:** A *generic* condition on  $\mathbf{X}$  depending on a parameter  $\mathbf{z} \in \Omega$  holds almost everywhere, i.e., if the condition doesn't hold for  $\mathbf{z} \in \Sigma \subset \Omega$ , then  $\mu(\Sigma) = 0$  with  $\mu$  a measure absolute continuous w.r.t. the Lebesgue one.

## Deterministic conditions

- Statistical independence → Independent Component Analysis

$$\mathbf{X} = \mathbf{M} \mathbf{S}_{\text{Ind}}^T$$

- Nonnegativity → Nonnegative Matrix Factorization

$$\mathbf{X} = \mathbf{M}_+ \mathbf{S}_+^T$$

- Sparsity → Sparse Component Analysis

$$\mathbf{X} = \mathbf{M} \mathbf{S}_{\text{Max0}}^T$$

- ...

## General case

$$\mathbf{X} = \mathbf{M} \mathbf{S}^T$$

---

$$\underbrace{\mathbf{X}}_{\mathbf{X}} = \underbrace{\mathbf{M} \mathbf{A}}_{\mathbf{M}_1} \underbrace{\mathbf{A}^{-1} \mathbf{S}^T}_{\mathbf{S}_1^T}$$

Uniqueness isn't guaranteed!

## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

Algorithm improvements

Exterior algebra

Numerical results



## Constraints for uniqueness

**Definition:** A *deterministic* condition on  $\mathbf{X}$  imposes a particular property of  $\mathbf{X}$  that is always true.

---

**Definition:** A *generic* condition on  $\mathbf{X}$  depending on a parameter  $\mathbf{z} \in \Omega$  holds almost everywhere, i.e., if the condition doesn't hold for  $\mathbf{z} \in \Sigma \subset \Omega$ , then  $\mu(\Sigma) = 0$  with  $\mu$  a measure absolute continuous w.r.t. the Lebesgue one.

## Problem statement

$$\mathbf{X} = \begin{bmatrix} | & \text{---} \\ & \mathbf{s}_1 \\ | & \text{---} \\ \mathbf{m}_1 & \end{bmatrix} + \dots + \begin{bmatrix} | & \text{---} \\ & \mathbf{s}_R \\ | & \text{---} \\ \mathbf{m}_R & \end{bmatrix}$$

---

$$\mathbf{X} = \mathbf{M}(\mathbf{z})\mathbf{S}^T(\mathbf{z}) = \sum_{r=1}^R \mathbf{m}_r(\mathbf{z}) \otimes \mathbf{s}(\boldsymbol{\xi}_r)$$

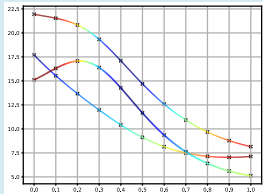
where

- $\mathbf{z} \in \Omega$  a subset of  $\mathbb{R}^n$
- $\mathbf{m}_r(\mathbf{z})$  are linearly independent
- each  $\mathbf{s}_r(\mathbf{z})$  depends on  $\ell$  independent parameters, entries of  $\boldsymbol{\xi}_r \in \mathbb{R}^\ell$
- each  $\mathbf{s}_r(\mathbf{z}) = \mathbf{s}_r(\boldsymbol{\xi}_r)$  has the structure

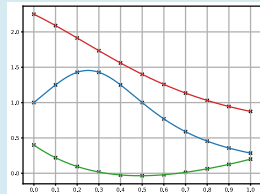
$$\mathbf{s}(\boldsymbol{\xi}_r) = \left[ \frac{p_1}{q_1} \circ \mathbf{f}(\boldsymbol{\xi}_r) \quad \dots \quad \frac{p_N}{q_N} \circ \mathbf{f}(\boldsymbol{\xi}_r) \right]$$

with  $p_h, q_h$  polynomials and  $\mathbf{f} = [f_1, \dots, f_\ell]$  a vectorial function.

# Example



$$= \begin{bmatrix} 8 & 1 & 6 \\ 3 & 5 & 7 \\ 4 & 9 & 2 \end{bmatrix} \times$$



Given the observed mixtures, we assume that

- the mixture matrix  $\mathbf{M}$  is constant and full rank;
- the source signals can be modeled by rational functions, i.e., the columns of  $\mathbf{S}$  are sampled of

$$s(t) = \frac{a_0 + a_1 t + \dots + a_p t^p}{b_0 + b_1 t + \dots + b_q t^q} \quad \text{with} \quad a_i, b_i \in \mathbb{R}, \quad t \in [t_b, t_e]$$



■  $\xi = [a_0, \dots, a_p, b_0, \dots, b_q]$

■  $\ell = p + q + 2$

■  $\mathbf{f}$  is the identity

Let  $\mathbf{t}(\mathbf{x}) = \left[ \frac{p_1}{q_1}(\mathbf{x}) \quad \dots \quad \frac{p_N}{q_N}(\mathbf{x}) \right]^T$  for  $\mathbf{x} \in \Theta = \{\mathbf{x} \in \mathbb{C}^\ell : q_1(\mathbf{x}) \cdots q_N(\mathbf{x}) \neq 0\}$ , if

1.  $\text{rank} \mathbf{M}(\mathbf{z}) = R$  for a generic choice of  $\mathbf{z}$
2. each  $f_h$  is the ratio of two analytical functions on  $\mathbb{C}^\ell$
3. there exists  $\boldsymbol{\xi}_0 \in \mathbb{C}^\ell$  s.t.  $\det \mathbf{J}(\mathbf{f}, \boldsymbol{\xi}_0) \neq 0$
4. the dimension of the span of  $\mathbf{t}(\mathbf{x})$  for  $\mathbf{x} \in \Theta$  is at least  $\hat{N}$
5.  $\text{rank} \mathbf{J}(\mathbf{t}, \mathbf{x}) > \hat{\ell}$  for a generic choice of  $\mathbf{z}$
6.  $R \leq \hat{N} - \hat{\ell}$

then

$$\mathbf{X} = \sum_{r=1}^R \mathbf{m}_r(\mathbf{z}) \otimes \mathbf{s}(\boldsymbol{\xi}_r)$$

is generically unique.

## Remarks for BSS

It is assumed that the columns of  $\mathbf{S}$  are values of the rational function

$$\mathbf{t} : \mathbf{x} \rightarrow \left[ \frac{p_1}{q_1}(\mathbf{x}) \quad \dots \quad \frac{p_N}{q_N}(\mathbf{x}) \right]^T.$$



The columns of  $\mathbf{S}$  belong to an algebraic variety  $\mathcal{V}$  which is described by a finite system of polynomials  $\{P_k\}_{k=1}^K$

$$\mathcal{V} = \left\{ (z_1, \dots, z_N) \in \mathbb{C}^N : P_k(z_1, \dots, z_N) = 0 \right\}.$$

## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

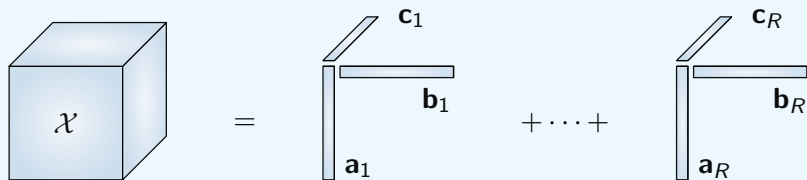
Algorithm improvements

Exterior algebra

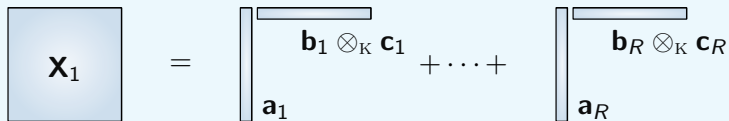
Numerical results

## Link with the CPD

$$\mathcal{X} = \mathbf{a}_1 \otimes \mathbf{b}_1 \otimes \mathbf{c}_1 + \dots + \mathbf{a}_R \otimes \mathbf{b}_R \otimes \mathbf{c}_R$$



$$\mathbf{X}_1 = \mathbf{a}_1 \otimes (\mathbf{b}_1 \otimes_{\mathbf{K}} \mathbf{c}_1)^T + \dots + \mathbf{a}_R \otimes (\mathbf{b}_R \otimes_{\mathbf{K}} \mathbf{c}_R)^T$$



$$(\mathbf{b}_r \otimes_{\mathbf{K}} \mathbf{c}_r) \in \mathcal{V} = \left\{ \text{vec}(\mathbf{Z}) : \begin{vmatrix} z_{i_1 j_1} & z_{i_1 j_2} \\ z_{i_2 j_1} & z_{i_2 j_2} \end{vmatrix} = 0 \right\}$$

## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

Algorithm improvements

Exterior algebra

Numerical results



## Algebraic algorithm outline

$$\underbrace{\mathbf{X}}_{\text{Known}} = \underbrace{\mathbf{M}}_{\text{Unknown}} \underbrace{\mathbf{S}^T}_{\text{Unknown}}$$

1. compute  $\mathbf{M}^{-1}$  from  $\mathbf{X}$ ;
2. compute  $\mathbf{S}$  as  $\mathbf{M}^{-1}\mathbf{X}$  transposed

## Equivalent condition I

$\mathbf{a}$  is a column of  $\mathbf{M}^{-1}$  if and only if  $\mathbf{X}^T \mathbf{a}$  is equal to a column of  $\mathbf{S}$



$$\mathbf{X}^T \mathbf{a} = (\mathbf{x}_1^T \mathbf{a}, \dots, \mathbf{x}_N^T \mathbf{a}) = (z_1, \dots, z_N) \in \mathcal{V}$$



$$P_k(\mathbf{x}_1^T \mathbf{a}, \dots, \mathbf{x}_N^T \mathbf{a}) = 0 \text{ for } k = 1, \dots, K$$



$$P_k^\otimes(\mathbf{x}_1^T, \dots, \mathbf{x}_N^T)(\mathbf{a} \otimes \dots \otimes \mathbf{a}) = 0 \text{ for } k = 1, \dots, K$$

where  $P_k^\otimes(\mathbf{x}_1^T, \dots, \mathbf{x}_N^T)$  is the vector obtained by formal substitution of  $(z_1, \dots, z_N)$  by  $\mathbf{x}_1^T, \dots, \mathbf{x}_N^T$  and the scalar multiplication by the tensor product.

## Equivalent condition II

$\mathbf{a}$  is a column of  $\mathbf{M}^{-1}$  if and only if  $\mathbf{X}^T \mathbf{a}$  is equal to a column of  $\mathbf{S}$



$$P_k^\otimes(\mathbf{x}_1^T, \dots, \mathbf{x}_N^T)(\mathbf{a} \otimes \dots \otimes \mathbf{a}) = 0 \text{ for } k = 1, \dots, K$$



$$\mathbf{Q} \text{vec}(\mathbf{a}^{\otimes p}) = \begin{bmatrix} P_1^\otimes(\mathbf{x}_1^T, \dots, \mathbf{x}_N^T) \\ \vdots \\ P_K^\otimes(\mathbf{x}_1^T, \dots, \mathbf{x}_N^T) \end{bmatrix} \text{vec}(\mathbf{a} \otimes \dots \otimes \mathbf{a}) = 0$$



The columns of  $\mathbf{M}^{-1}$  belong to the intersection of  $\mathbf{Q}$  kernel and the subspace of vectorized order  $p$  symmetric tensors.

The diagram illustrates the equation  $\mathbf{X} = \mathbf{M} \mathbf{S}^T$ . Each term is enclosed in a light blue rectangular box. Below the box for  $\mathbf{X}$  is a bracket labeled "Known" in green text. Below the box for  $\mathbf{M}$  is a bracket labeled "Unknown" in red text. Below the box for  $\mathbf{S}^T$  is a bracket labeled "Unknown" in red text. An equals sign is positioned between the boxes for  $\mathbf{X}$  and  $\mathbf{M}$ .

1. compute  $\mathbf{M}^{-1}$  from  $\mathbf{X}$ ;
  - 1.1 compute  $\mathbf{Q}$ ;
  - 1.2 compute  $\text{null}(\mathbf{Q})$  intersected with the space of vectorized symmetric tensors;
2. compute  $\mathbf{S}$  as  $\mathbf{M}^{-1}\mathbf{X}$  transposed.

## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

**The bottleneck**

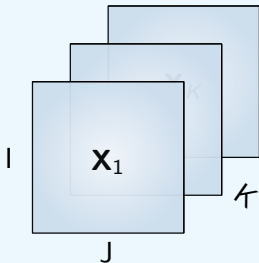
Algorithm improvements

Exterior algebra

Numerical results

## Q for CPD

Let  $\mathcal{X}$  be an order-3 tensor of dimension  $(I \times J \times K)$ , then



**Definition:**  $\mathbf{Q}$  is a  $C_I^2 C_J^2 \times C_{K+1}^2$  matrix whose  $k$ -th column can be written as

$$\mathbf{Q}(\cdot, k) = \text{vec}(\mathcal{C}_2(\mathbf{X}_{k_1} + \mathbf{X}_{k_2}) - \mathcal{C}_2(\mathbf{X}_{k_1}) - \mathcal{C}_2(\mathbf{X}_{k_2}))$$

where

- $(k_1, k_2)$  is the  $k$ -th element of  $\mathcal{Q}_K^2 = \{(k_1, k_2) : 1 \leq k_1 \leq k_2 \leq K\}$ ;
- $C_N^2$  is the binomial of  $N$  over 2;
- $\mathcal{C}_2(\mathbf{X}_h)$  is the matrix with the determinants of every  $2 \times 2$  minors of  $\mathbf{X}_h$ .

## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

Algorithm improvements

Exterior algebra

Numerical results

## Example

Let  $\mathcal{X}$  be an order-3 tensor of dimension  $(I \times 3 \times 2)$  such that

$$\mathbf{X}_1 = [\mathbf{a}_1 \quad \mathbf{a}_2 \quad \mathbf{a}_3] \quad \text{and} \quad \mathbf{X}_2 = [\mathbf{b}_1 \quad \mathbf{b}_2 \quad \mathbf{b}_3]$$

The matrix  $\mathbf{Q} = [\mathbf{q}_1 \quad \mathbf{q}_2 \quad \mathbf{q}_3]$  associated with  $\mathcal{X}$  slices is such that

- $\mathbf{q}_1 = \text{vec}(\mathcal{C}_2(\mathbf{X}_1 + \mathbf{X}_1) - \mathcal{C}_2(\mathbf{X}_1) - \mathcal{C}_2(\mathbf{X}_1)) = 2\text{vec}(\mathcal{C}_2(\mathbf{X}_1))$
- $\mathbf{q}_2 = \text{vec}(\mathcal{C}_2(\mathbf{X}_1 + \mathbf{X}_2) - \mathcal{C}_2(\mathbf{X}_1) - \mathcal{C}_2(\mathbf{X}_2))$
- $\mathbf{q}_3 = \text{vec}(\mathcal{C}_2(\mathbf{X}_2 + \mathbf{X}_2) - \mathcal{C}_2(\mathbf{X}_2) - \mathcal{C}_2(\mathbf{X}_2)) = 2\text{vec}(\mathcal{C}_2(\mathbf{X}_2))$

The compound matrices are

- $\mathcal{C}_2(\mathbf{X}_1) = [\alpha_1 \quad \alpha_2 \quad \alpha_3]$
- $\mathcal{C}_2(\mathbf{X}_2) = [\beta_1 \quad \beta_2 \quad \beta_3]$
- $\mathcal{C}_2(\mathbf{X}_1 + \mathbf{X}_2) = [\gamma_1 \quad \gamma_2 \quad \gamma_3]$

$$\Rightarrow \mathbf{Q} = \begin{bmatrix} 2\alpha_1 & \gamma_1 - (\alpha_1 + \beta_1) & 2\beta_1 \\ 2\alpha_2 & \gamma_2 - (\alpha_2 + \beta_2) & 2\beta_2 \\ 2\alpha_3 & \gamma_3 - (\alpha_3 + \beta_3) & 2\beta_3 \end{bmatrix}$$



## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

Algorithm improvements

Exterior algebra

Numerical results

## Exterior algebra and $\mathcal{C}_2$

**Definition:** The exterior product  $\wedge : \mathbb{R}^I \times \mathbb{R}^J \rightarrow \mathbb{R}^{I \times J}$  is defined as

$$\mathbf{a} \wedge \mathbf{b} = \mathbf{a}\mathbf{b}^T - \mathbf{b}\mathbf{a}^T.$$

### Example

Let  $\mathbf{a}^T = [a \ b \ c]$  and  $\mathbf{b}^T = [d \ e \ f]$  be two vectors, then

$$\mathbf{a} \wedge \mathbf{b} = \begin{bmatrix} 0 & ae - bd & af - cd \\ -ae + bd & 0 & bf - ce \\ -af + cd & -bf + ce & 0 \end{bmatrix}$$

The matrix  $\mathbf{a} \wedge \mathbf{b}$  has the properties:

- the diagonal entries are zeros;
- it is skew-symmetric;
- the elements highlighted are the entries of  $\mathcal{C}_2([\mathbf{a} \ \mathbf{b}])$ ;

## Algebraic algorithm optimization steps

**Property:**  $\langle \text{vec}(\mathcal{C}_2([\mathbf{a} \ \mathbf{b}])), \text{vec}(\mathcal{C}_2([\mathbf{a} \ \mathbf{b}])) \rangle = 2 \langle \text{vec}(\mathbf{a} \wedge \mathbf{b}), \text{vec}(\mathbf{a} \wedge \mathbf{b}) \rangle$   
 $= 4\|\mathbf{a}\|^2\|\mathbf{b}\|^2 - \langle \mathbf{a}, \mathbf{b} \rangle^2$

---

**Idea:** Using the exterior product to improve the algebraic algorithm.

- $\mathbf{Q}$  is not explicitly constructed
- $\text{null}(\mathbf{Q}) = \text{null}(\mathbf{Q}^T \mathbf{Q})$
- $\langle \mathbf{q}_k, \mathbf{q}_h \rangle = \langle \text{vec}(\mathcal{C}_2(\mathbf{Y})), \text{vec}(\mathcal{C}_2(\mathbf{Z})) \rangle$  with  $\mathbf{Y}, \mathbf{Z}$  slices or a linear combination of slices of the input tensor  $\mathcal{X}$



- Pre-compute  $\mathbf{A}_{kh} = \mathbf{X}_k^T \mathbf{X}_h$  for  $h \leq k$
- define 6 functions  $f_i$  that depends on the  $\mathbf{A}_{kh}$
- use  $f_i$  and  $\mathbf{A}_{kh}$  to compute the entries of  $\mathbf{Q}^T \mathbf{Q}$

## Overview

The Blind Source Separation

Deterministic uniqueness

Generic uniqueness

The Canonical Polyadic Decomposition

From the theorem to the algorithm

The bottleneck

Algorithm improvements

Exterior algebra

**Numerical results**

## Preliminary numerical results (work in progress)

Given an input tensor  $\mathcal{X}$  of dimensions  $(I \times J \times K)$  such that its rank is  $I = (J - 1)(K - 1)$

<b>dimensions</b>	<b>rank</b>	<b>CPU time</b>
$20 \times 5 \times 6$	$R = 20$	5.16sec
$24 \times 5 \times 7$	$R = 24$	24.42sec
$25 \times 6 \times 6$	$R = 25$	21.82sec
$28 \times 5 \times 8$	$R = 28$	283.22sec
$30 \times 6 \times 7$	$R = 30$	$\approx 7\text{min}$




The link between rank and dimensions is meant to satisfy [Domanov and De Lathauwer 2016] theorem.

## Conclusive remarks

- conditions that guarantee generic uniqueness;
- from the theorem structure to a CPD algorithm;
- bottleneck due to compound matrices;
- optimization based on the exterior product.

Thank you for the attention!  
Questions?

## References I

-  Comon, P. and C. Jutten (2009). *Handbook of blind source separation: Independent component analysis and applications*. Academic press.
-  Domanov, I. and L. De Lathauwer (July 2013). “On the uniqueness of the canonical polyadic decomposition of third-order tensors — Part I: Basic results and uniqueness of one factor matrix”. In: *SIAM J. Matrix Anal. Appl.* 34.3, pp. 855–875.
-  — (June 2016). “Generic Uniqueness of a Structured Matrix Factorization and Applications in Blind Source Separation”. In: *IEEE J. Sel. Topics Signal Process.* 10.4, pp. 701–711.