

A subspace-conjugate gradient method for Kronecker-structured equations

Martina Iannacito

joint work with Davide Palitta and Valeria Simoncini

HAW visiting

January 28, 2026



Multiterm matrix equation – I

Consider the **multiterm Sylvester equation**

$$\mathbf{A}_1 \mathbf{X} \mathbf{B}_1 + \cdots + \mathbf{A}_m \mathbf{X} \mathbf{B}_m = \mathbf{C}$$

where $\mathbf{A}_k \in \mathbb{R}^{n_A \times n_A}$ and $\mathbf{B}_k \in \mathbb{R}^{n_B \times n_B}$ are symmetric, and \mathbf{C} is low-rank (s_C).

Consider the **multiterm Sylvester equation**

$$\mathbf{A}_1\mathbf{X}\mathbf{B}_1 + \cdots + \mathbf{A}_m\mathbf{X}\mathbf{B}_m = \mathbf{C}$$

where $\mathbf{A}_k \in \mathbb{R}^{n_A \times n_A}$ and $\mathbf{B}_k \in \mathbb{R}^{n_B \times n_B}$ are symmetric, and \mathbf{C} is low-rank (s_C).

They arise in different fields, e.g.,

- statistics and probability
- control and system theory
- signal processing
- uncertainty quantification
- physical phenomena simulation

Multiterm matrix equation in control and system theory

Consider the continuous-time linear system

$$\begin{aligned}\dot{\mathbf{x}} &= \mathbf{A}\mathbf{x} + \mathbf{B}_1\mathbf{u} \\ \mathbf{y} &= \mathbf{B}_2^*\mathbf{x}\end{aligned}\tag{1}$$

where \mathbf{x} is the *model state*, \mathbf{u} is the *input*, \mathbf{y} is the *output*, \mathbf{A} , \mathbf{B}_1 and \mathbf{B}_2 are invariant. If \mathbf{A} is stable¹, then the solutions \mathbf{P} and \mathbf{Q} of the Lyapunov equations

$$\begin{aligned}\mathbf{A}\mathbf{P} + \mathbf{P}\mathbf{A}^* + \mathbf{B}_1\mathbf{B}_1^* &= 0 \\ \mathbf{A}^*\mathbf{Q} + \mathbf{Q}\mathbf{A} + \mathbf{B}_2\mathbf{B}_2^* &= 0\end{aligned}$$

are called *controllability* and *observability* Gramians, respectively. They can be used to measure the energy transfers in the system (1).

See [Antoulas 2005; Simoncini 2016] for further details.

Solving multiterm matrix equations

If the number of terms $m = 2$

- if n_A and n_B are small, direct methods can be used, e.g., Bartels–Stewart method
- otherwise, need iterative methods, e.g., projection methods

Solving multiterm matrix equations

If the number of terms $m = 2$

- if n_A and n_B are small, direct methods can be used, e.g., Bartels–Stewart method
- otherwise, need iterative methods, e.g., projection methods

If the number of terms $m > 2$

- direct methods can be used only through the Kronecker form, that is

$$(\mathbf{B}_1^\top \otimes_{\mathbf{K}} \mathbf{A}_1 + \cdots + \mathbf{B}_m^\top \otimes_{\mathbf{K}} \mathbf{A}_m) \text{vec}(\mathbf{X}) = \text{vec}(\mathbf{C}).$$

prohibitive for moderate n_A, n_B .

Solving multiterm matrix equations

If the number of terms $m = 2$

- if n_A and n_B are small, direct methods can be used, e.g., Bartels–Stewart method
- otherwise, need iterative methods, e.g., projection methods

If the number of terms $m > 2$

- direct methods can be used only through the Kronecker form, that is

$$(\mathbf{B}_1^\top \otimes_{\mathbb{K}} \mathbf{A}_1 + \cdots + \mathbf{B}_m^\top \otimes_{\mathbb{K}} \mathbf{A}_m) \text{vec}(\mathbf{X}) = \text{vec}(\mathbf{C}).$$

prohibitive for moderate n_A, n_B .

- iterative methods, such as
 - matrix-oriented Krylov methods
 - projection methods
 - fixed-point iterations
 - Riemannian optimization schemes

Let $\mathcal{L}(\mathbf{X}) = \mathbf{A}_1\mathbf{X}\mathbf{B}_1 + \cdots + \mathbf{A}_m\mathbf{X}\mathbf{B}_m$ for any $\mathbf{X} \in \mathbb{R}^{n \times n}$, with $n = n_A = n_B$, and $\mathbf{A}_h, \mathbf{B}_h$ are symmetric.

Let $\mathcal{L}(\mathbf{X}) = \mathbf{A}_1\mathbf{X}\mathbf{B}_1 + \cdots + \mathbf{A}_m\mathbf{X}\mathbf{B}_m$ for any $\mathbf{X} \in \mathbb{R}^{n \times n}$, with $n = n_A = n_B$, and $\mathbf{A}_h, \mathbf{B}_h$ are symmetric.

Let $(\mathcal{L}(\mathbf{Y}))^\top = \mathcal{L}(\mathbf{Y})$ for \mathbf{Y} symmetric, so that we can work with $\mathbf{Y} = \mathbf{Y}\mathbf{Y}^\top$ with $\mathbf{Y} \in \mathbb{R}^{n \times r}$.

Matrix-oriented CG – I

Let $\mathcal{L}(\mathbf{X}) = \mathbf{A}_1\mathbf{X}\mathbf{B}_1 + \cdots + \mathbf{A}_m\mathbf{X}\mathbf{B}_m$ for any $\mathbf{X} \in \mathbb{R}^{n \times n}$, with $n = n_A = n_B$, and $\mathbf{A}_h, \mathbf{B}_h$ are symmetric.

Let $(\mathcal{L}(\mathbf{Y}))^\top = \mathcal{L}(\mathbf{Y})$ for \mathbf{Y} symmetric, so that we can work with $\mathbf{Y} = \mathbf{Y}\mathbf{Y}^\top$ with $\mathbf{Y} \in \mathbb{R}^{n \times r}$.

Given an initial guess $\mathbf{X}_0 \in \mathbb{R}^{n \times n}$, the matrix-oriented CG iterate is

$$\mathbf{X}_{k+1} = \mathbf{X}_k + \alpha_k \mathbf{P}_k$$

where $\alpha_k \in \mathbb{R}$.

The direction matrix iterate is

$$\mathbf{P}_{k+1} = \mathbf{R}_{k+1} + \beta_k \mathbf{P}_k$$

where $\mathbf{R}_{k+1} = \mathbf{C} - \mathcal{L}(\mathbf{X}_{k+1})$ is the residual and $\beta_k \in \mathbb{R}$.

Matrix-oriented CG – II

The dense matrices

- \mathbf{X}_{k+1} for the *iterative solution*
- \mathbf{R}_{k+1} for the *residual*
- \mathbf{P}_{k+1} for the *directions*

have size $(n \times n)$ matrices with n potentially very large



If \mathbf{C} is low-rank, $\mathbf{C} = \mathbf{C}\mathbf{C}^\top$ with \mathbf{C} an $(n \times s)$ matrix, we can use the *factored forms*

$$\mathbf{X}_{k+1} = [\mathbf{X}_k, \sqrt{\alpha_k} \mathbf{P}_k][\mathbf{X}_k, \sqrt{\alpha_k} \mathbf{P}_k]^\top$$

where $\mathbf{X}_k = \mathbf{X}_k \mathbf{X}_k^\top$ and $\mathbf{P}_k = \mathbf{P}_k \mathbf{P}_k^\top$ with \mathbf{X}_k and \mathbf{P}_k thin matrices.

Usually the blocks X_{k+1} , R_{k+1} and P_{k+1} get larger accumulating redundant information

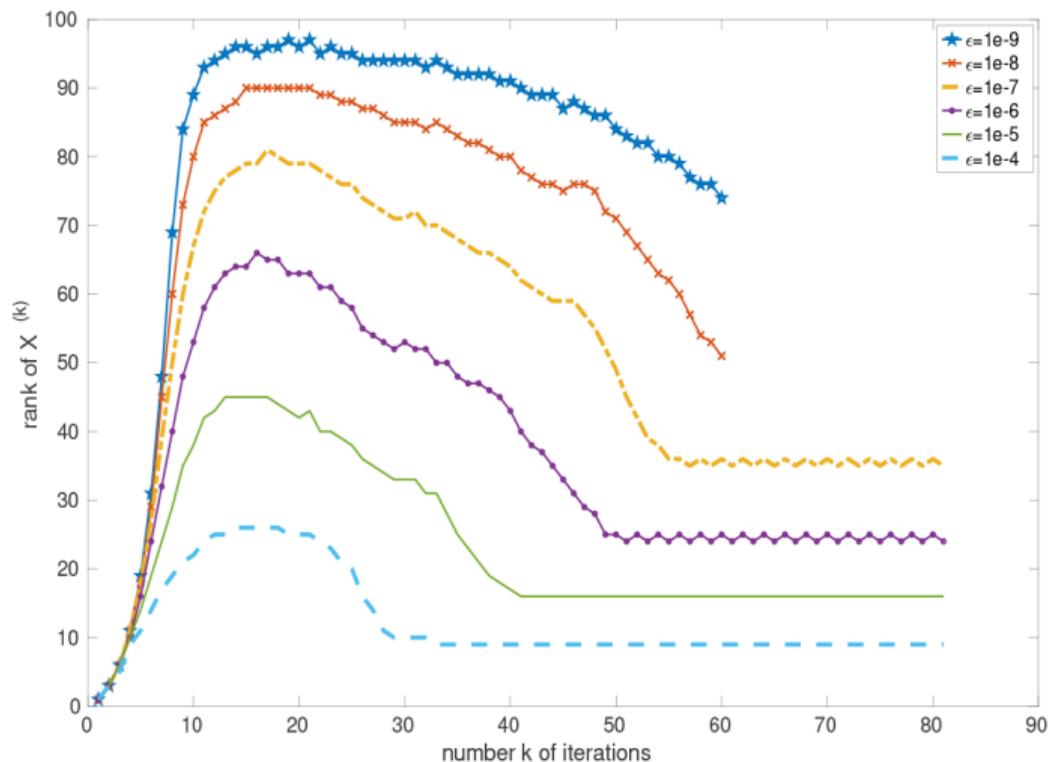


Introduce truncation on X_{k+1} , R_{k+1} and P_{k+1} at fixed rank or precision

While computational costs are controlled, side-effects are

- delayed or stagnating convergence;
- hard to control the rank.

Rank vs iteration [Simoncini and Hao 2023, Fig. 4]



Subspace CG

Idea: modifying α_k and β_k nature to improve convergence

Subspace CG

Idea: modifying α_k and β_k nature to improve convergence

Define $\Phi : \mathbb{R}^{n \times n} \rightarrow \mathbb{R}$ such that

$$\Phi(\mathbf{X}) = \frac{1}{2} \langle \mathcal{L}(\mathbf{X}), \mathbf{X} \rangle - \langle \mathbf{C}, \mathbf{X} \rangle.$$

Consider the minimization problem

$$\text{Find } \mathbf{X} \in \mathbb{R}^{n \times n} \text{ s.t. } \mathbf{X} = \arg \min_{\mathbf{X} \in \mathbb{R}^{n \times n}} |\Phi(\mathbf{X})|.$$

and the new recurrence

$$\mathbf{X}_{k+1} = \mathbf{X}_k + P_k \alpha_k P_k^\top$$

where $\alpha_k \in \mathbb{R}^{s_k \times s_k}$ and $P_k \in \mathbb{R}^{n \times s_k}$.

- recurrence for the direction

$$\mathbf{P}_{k+1} = \mathbf{R}_{k+1} + P_k \beta_k P_k^\top$$

where the residual remains $\mathbf{R}_{k+1} = \mathbf{C} - \mathcal{L}(\mathbf{X}_{k+1})$ and $\beta_k \in \mathbb{R}^{s_k \times s_k}$.

- the direction factor matrix $P_{k+1} \in \mathbb{R}^{n \times s_{k+1}}$ is computed from

$$\mathbf{P}_{k+1} = P_{k+1} \gamma_{k+1} P_{k+1}^\top$$

where $\gamma_{k+1} \in \mathbb{R}^{s_{k+1} \times s_{k+1}}$.

Why matrix coefficients (might) work better?

Let $\mathbf{x}_j^{(k+1)}$ denotes the j -th column of \mathbf{X}_{k+1} , the iterative solution at the $k + 1$ iteration.

In the matrix-CG, $\mathbf{x}_j^{(k+1)} = \mathbf{x}_j^{(k)} + \mathbf{u}$ where \mathbf{u} is the updated defined as

$$\mathbf{u} = \alpha_k (p_{j,1} \mathbf{p}_1 + \cdots + p_{j,s_k} \mathbf{p}_{s_k})$$

where $P_k = [\mathbf{p}_1, \dots, \mathbf{p}_{s_k}]$ and $P_k(i, j) = p_{i,j}$.

Why matrix coefficients (might) work better?

Let $\mathbf{x}_j^{(k+1)}$ denotes the j -th column of \mathbf{X}_{k+1} , the iterative solution at the $k + 1$ iteration.

In the matrix-CG, $\mathbf{x}_j^{(k+1)} = \mathbf{x}_j^{(k)} + \mathbf{u}$ where \mathbf{u} is the updated defined as

$$\mathbf{u} = \alpha_k (p_{j,1} \mathbf{p}_1 + \cdots + p_{j,s_k} \mathbf{p}_{s_k})$$

where $P_k = [\mathbf{p}_1, \dots, \mathbf{p}_{s_k}]$ and $P_k(i, j) = p_{i,j}$.

In the new method, the j -th column of \mathbf{X}_{k+1} is computed as $\mathbf{x}_j^{(k+1)} = \mathbf{x}_j^{(k)} + \mathbf{u}$ where the updated \mathbf{u} is

$$\left(\sum_{i=1}^{s_k} \alpha_{1,i} p_{j,i} \right) \mathbf{p}_1 + \cdots + \left(\sum_{i=1}^{s_k} \alpha_{s_k,i} p_{j,i} \right) \mathbf{p}_{s_k}$$

with $\alpha_k(i, j) = \alpha_{i,j}$.

New matrix coefficients – I

- α_k is the solution of

$$\min_{\alpha \in \mathbb{R}^{s_k \times s_k}} \Phi(\mathbf{X}_k + P_k \alpha P_k^\top)$$

that corresponds the orthogonality condition

$$\text{vec}(\mathbf{R}_{k+1}) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k).$$

New matrix coefficients – I

- α_k is the solution of

$$\min_{\alpha \in \mathbb{R}^{s_k \times s_k}} \Phi(\mathbf{X}_k + P_k \alpha P_k^\top)$$

that corresponds the orthogonality condition

$$\text{vec}(\mathbf{R}_{k+1}) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k).$$

In the matrix-CG, $\alpha_k \in \mathbb{R}$ is computed to satisfy the orthogonality condition

$$\text{vec}(\mathbf{R}_k - \alpha_k \mathcal{L}(\mathbf{P}_k)) \perp \text{vec}(\mathbf{P}_k).$$

recalling that $\mathbf{R}_{k+1} = \mathbf{R}_k - \alpha_k \mathcal{L}(\mathbf{P}_k)$ and $\mathbf{P}_k = P_k P_k^\top$.

New matrix coefficients – I

- α_k is the solution of

$$\min_{\alpha \in \mathbb{R}^{s_k \times s_k}} \Phi(\mathbf{X}_k + P_k \alpha P_k^\top)$$

that corresponds the orthogonality condition

$$\text{vec}(\mathbf{R}_{k+1}) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k). \quad \leftarrow \text{orthogonality w.r.t. subspace of size } s_k^2 \text{ of } \mathbb{R}^{n^2}$$

In the matrix-CG, $\alpha_k \in \mathbb{R}$ is computed to satisfy the orthogonality condition

$$\text{vec}(\mathbf{R}_k - \alpha_k \mathcal{L}(\mathbf{P}_k)) \perp \text{vec}(\mathbf{P}_k).$$

recalling that $\mathbf{R}_{k+1} = \mathbf{R}_k - \alpha_k \mathcal{L}(\mathbf{P}_k)$ and $\mathbf{P}_k = P_k P_k^\top$.

New matrix coefficients – I

- α_k is the solution of

$$\min_{\alpha \in \mathbb{R}^{s_k \times s_k}} \Phi(\mathbf{X}_k + P_k \alpha P_k^\top)$$

that corresponds the orthogonality condition

$$\text{vec}(\mathbf{R}_{k+1}) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k). \quad \leftarrow \text{orthogonality w.r.t. subspace of size } s_k^2 \text{ of } \mathbb{R}^{n^2}$$

In the matrix-CG, $\alpha_k \in \mathbb{R}$ is computed to satisfy the orthogonality condition

$$\text{vec}(\mathbf{R}_k - \alpha_k \mathcal{L}(\mathbf{P}_k)) \perp \text{vec}(\mathbf{P}_k). \quad \leftarrow \text{orthogonality w.r.t. subspace of size } 1 \text{ of } \mathbb{R}^{n^2}$$

recalling that $\mathbf{R}_{k+1} = \mathbf{R}_k - \alpha_k \mathcal{L}(\mathbf{P}_k)$ and $\mathbf{P}_k = P_k P_k^\top$.

- β_k is s.t. the updated directions, \mathbf{P}_{k+1} , are \mathcal{L} -orthogonal w.r.t. the previous ones,

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k)$$

- β_k is s.t. the updated directions, \mathbf{P}_{k+1} , are \mathcal{L} -orthogonal w.r.t. the previous ones,

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k)$$

In the matrix-CG, $\beta_k \in \mathbb{R}$ is computed to satisfy the orthogonality condition

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{vec}(\mathbf{P}_k).$$

- β_k is s.t. the updated directions, \mathbf{P}_{k+1} , are \mathcal{L} -orthogonal w.r.t. the previous ones,

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{range}(P_k \otimes_K P_k) \quad \leftarrow \text{orthogonality w.r.t. subspace of size } s_k^2 \text{ of } \mathbb{R}^{n^2}$$

In the matrix-CG, $\beta_k \in \mathbb{R}$ is computed to satisfy the orthogonality condition

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{vec}(\mathbf{P}_k).$$

- β_k is s.t. the updated directions, \mathbf{P}_{k+1} , are \mathcal{L} -orthogonal w.r.t. the previous ones,

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{range}(P_k \otimes_{\mathbb{K}} P_k) \quad \leftarrow \text{orthogonality w.r.t. subspace of size } s_k^2 \text{ of } \mathbb{R}^{n^2}$$

In the matrix-CG, $\beta_k \in \mathbb{R}$ is computed to satisfy the orthogonality condition

$$\text{vec}(\mathcal{L}(\mathbf{P}_{k+1})) \perp \text{vec}(\mathbf{P}_k). \quad \leftarrow \text{orthogonality w.r.t. subspace of size } 1 \text{ of } \mathbb{R}^{n^2}$$

Computation of the matrix coefficients - I

- α_k is the unique solution of

$$P_k^\top \mathcal{L}(P_k \alpha P_k^\top - \mathbf{R}_k) P_k = 0$$

Computation of the matrix coefficients - I

- α_k is the unique solution of

$$P_k^\top \mathcal{L}(P_k \alpha P_k^\top - \mathbf{R}_k) P_k = 0$$

or equivalently

$$P_k^\top \mathcal{L}(P_k \alpha P_k^\top) P_k = P_k^\top \mathbf{R}_k P_k.$$

Computation of the matrix coefficients - I

- α_k is the unique solution of

$$P_k^\top \mathcal{L}(P_k \alpha P_k^\top - \mathbf{R}_k) P_k = 0$$

or equivalently

$$P_k^\top \mathcal{L}(P_k \alpha P_k^\top) P_k = P_k^\top \mathbf{R}_k P_k.$$

- β_k is the unique solution of

$$P_k^\top \mathcal{L}(P_k \beta P_k^\top + \mathbf{R}_{k+1}) P_k = 0$$

or equivalently

$$P_k^\top \mathcal{L}(P_k \beta P_k^\top) P_k = -P_k^\top \mathcal{L}(\mathbf{R}_{k+1}) P_k.$$

Computation of the matrix coefficients - II

Both equations defining α_k and β_k are linear matrix equations of the **same type** as the original problem but of **smaller size** ($s_k \times s_k$).

Let $\tilde{\mathbf{A}}_h^{(k)}$ and $\tilde{\mathbf{B}}_h^{(k)}$ be ($s_k \times s_k$) matrices for $h = 1, \dots, \ell$ defined as

$$\tilde{\mathbf{A}}_h^{(k)} = P_k^\top \mathbf{A}_h P_k \quad \text{and} \quad \tilde{\mathbf{B}}_h^{(k)} = P_k^\top \mathbf{B}_h P_k.$$

Then α_k and β_k are solutions of

$$\begin{aligned} \tilde{\mathbf{A}}_1^{(k)} \alpha \tilde{\mathbf{B}}_1^{(k)} + \dots + \tilde{\mathbf{A}}_m^{(k)} \alpha \tilde{\mathbf{B}}_m^{(k)} &= P_k^\top \mathbf{R}_k P_k \\ \tilde{\mathbf{A}}_1^{(k)} \beta \tilde{\mathbf{B}}_1^{(k)} + \dots + \tilde{\mathbf{A}}_m^{(k)} \beta \tilde{\mathbf{B}}_m^{(k)} &= -P_k^\top \mathcal{L}(\mathbf{R}_{k+1}) P_k \end{aligned}$$

Practical enhancements

The entire method works also if $\mathcal{L}(\mathbf{X}) \neq (\mathcal{L}(\mathbf{X}))^\top$, but we distinguish between *left* (ℓ) and *right* (r) factors, e.g.,

$$\mathbf{X}_{k+1} = \mathbf{X}_k + P_k^{(\ell)} \alpha_k (P_k^{(r)})^\top.$$

Practical enhancements

The entire method works also if $\mathcal{L}(\mathbf{X}) \neq (\mathcal{L}(\mathbf{X}))^\top$, but we distinguish between *left* (ℓ) and *right* (r) factors, e.g.,

$$\mathbf{X}_{k+1} = \mathbf{X}_k + P_k^{(\ell)} \alpha_k (P_k^{(r)})^\top.$$

To make Ss-CG method efficient, we consider some enhancements

- block representation;
- truncation at a prescribed `maxrank` value;
- inexact coefficients;
- stopping criterion.

Inexact coefficients

- if `maxrank` is sufficiently small, α_k and β_k can be computed solving

$$\left(\sum_{j=1}^m \tilde{\mathbf{B}}_j^\top \otimes_{\mathbb{K}} \tilde{\mathbf{A}}_j \right) \text{vec}(\alpha) = \text{vec}(P_k^\top \mathbf{R}_k P_k)$$

$$\left(\sum_{j=1}^m \tilde{\mathbf{B}}_j^\top \otimes_{\mathbb{K}} \tilde{\mathbf{A}}_j \right) \text{vec}(\beta) = -\text{vec}(P_k^\top \mathcal{L}(\mathbf{R}_{k+1}) P_k)$$

Inexact coefficients

- if `maxrank` is sufficiently small, α_k and β_k can be computed solving

$$\left(\sum_{j=1}^m \tilde{\mathbf{B}}_j^\top \otimes_{\mathbb{K}} \tilde{\mathbf{A}}_j \right) \text{vec}(\alpha) = \text{vec}(P_k^\top \mathbf{R}_k P_k)$$

$$\left(\sum_{j=1}^m \tilde{\mathbf{B}}_j^\top \otimes_{\mathbb{K}} \tilde{\mathbf{A}}_j \right) \text{vec}(\beta) = -\text{vec}(P_k^\top \mathcal{L}(\mathbf{R}_{k+1}) P_k)$$

- otherwise, use an iterative solver at prescribed precision getting

$$\tilde{\alpha}_k = \alpha_k + \epsilon_k \quad \text{and} \quad \tilde{\beta}_k = \beta_k + \psi_k$$

Warning

The orthogonality properties previously seen are lost. Due to the use of truncation, inexactness of the coefficient matrices has not a strong effect.

Stopping criterion

Since the computation of the exact residual might be expensive, we test if the relative difference between consecutive iterates is smaller than the prescribed accuracy `tol`, that is

$$\|\mathbf{x}_{k+1} - \mathbf{x}_k\|_F / \|\mathbf{x}_{k+1}\|_F \leq \text{tol}$$

The difference between two iterates can be efficiently computed using the block structure and the trace properties!

Warning

It might be sensitive to ill-conditioning of the problem, thus the true residual is computed explicitly at completion.

Memory requirements

- each iterate in block format requires at most

$$(n_A + n_B)\text{maxrank} + \text{maxrank}^2 \quad \text{units of memory}$$

- matrix coefficients requires s_k^2 reaching at most maxrank units of memory
- dynamic truncation of the residual requires

$$(m \cdot \text{maxrank} + s_C)(n_A + n_B) \quad \text{units of memory}$$

while randomized truncation requires at least

$$\text{maxrankR}(n_A + n_B) \quad \text{units of memory}$$

Memory costs are comparable to those of the Truncated CG

The multiterm matrix equations of the following numerical examples are solved

- Ss-CG
 - deter(ministic)
 - rand(omized)
- TPCG: Truncated matrix-oriented Preconditioned CG [Kressner et al. 2011; Simoncini and Hao 2023]
- R-NLCG: Riemannian, nonlinear CG [Bioli et al. 2025]
- MultiRB: projection method for finite element discretizations of differential equations with stochastic input [Powell et al. 2017]

We consider the discretization of a heat model problem in $(0, 1)^2$, and resulting in the equation

$$\mathbf{A}\mathbf{X} + \mathbf{X}\mathbf{A} + \mathbf{M}\mathbf{X}\mathbf{M}^\top = \mathbf{C}$$

where \mathbf{A} is the discretization of the 2D Laplace operator, and $\mathbf{M} = \mathbf{N}\mathbf{N}^\top$ that allocates the Robin conditions $\bar{\mathbf{n}} \cdot \nabla(x) = \delta u(x - 1)$ on one of the domain boundaries, while zero Dirichlet conditions are imposed on the rest of the boundary.

In our experiments we consider $\delta \in \{0.5, 0.9\}$. We precondition with $\mathcal{L}_0(\mathbf{X}) = \mathbf{A}\mathbf{X} + \mathbf{X}\mathbf{A}$ by running $t_{ADI} = 8$ LR-ADI iterations.

Numerical example I – results

Example	n_A	maxrank	SSS (iter/alloc/rank)	TPCG	SsCG determ.	SsCG rand'zed
$\delta = 0.5$	102400	20	6.15 (7/126/31)	61.35 (16)	– (100)	– (100)
		30		22.78 (4)	17.93 (3)	18.17 (3)
	250000	20		– (100)	– (100)	– (100)
		30		60.84 (4)	64.46 (4)	64.45 (4)
$\delta = 0.9$	102400	40	– (50/ /)	– (100)	– (100)	– (100)
		50		310.72 (26)	58.11 (5)	58.52 (5)
		–				
	250000	50		2401.90 (93)	– (100)	– (100)
		60		936.39 (30)	119.55 (4)	120.43 (4)
		–		– (50/ /)		

– no conv.

Table: Running time in seconds, and in parenthesis the number of iterations. Stopping tolerance $\text{tol} = 10^{-6}$. For SSS, number of iterations, the subspace total memory allocation for length n_A vectors and the solution rank are reported.

Numerical example II – [Bioli et al. 2025]

The stationary diffusion equation

$$-\nabla \cdot (\kappa \nabla u) = 0 \quad \text{in} \quad (0, 1) \times (0, 1)$$

with Dirichlet boundary conditions and semi-separable diffusion coefficient:

$$\kappa(x, y) = \sum_{j=0}^m \delta_j \kappa_{x,j}(x) \kappa_{y,j}(y) = 1 + \sum_{j=1}^{m_k-1} \frac{10^j}{j!} x^j y^j.$$

The resulting multiterm linear equation is

$$\sum_{j=1}^{m_k} \delta_j (\mathbf{A}_{j,x} \mathbf{X} \mathbf{D}_{j,y} + \mathbf{D}_{j,x} \mathbf{X} \mathbf{A}_{j,y}) = \mathbf{C}$$

with \mathbf{C} being rank 4, $m_k = 4$ and a total of 8 terms.

Numerical example II – results

n type	Precond	maxrank	R-NLCG	TPCG	SsCG determ.	SsCG rand.
10000	\mathcal{P}_1	20	– (100)	– (100)	– (100)	– (100)
	\mathcal{P}_1	40	– (100)	– (100)	1.08 (5)	0.92 (5)
	\mathcal{P}_1	60	– (100)	– (100)	2.47 (5)	2.34 (5)
	\mathcal{P}_2	20	11.25 (36)	11.42 (38)	– (100)	– (100)
	\mathcal{P}_2	40	*42.97 (36)	15.54 (33)	– (100)	– (100)
	\mathcal{P}_2	60	*98.62 (35)	32.39 (28)	9.59 (5)	8.37 (5)
102400	\mathcal{P}_1	20	– (100)	– (100)	– (100)	– (100)
	\mathcal{P}_1	40	†	– (100)	18.17 (6)	8.74 (6)
	\mathcal{P}_1	60	†	– (100)	23.50 (5)	16.93 (5)
	\mathcal{P}_2	20	183.44 (41)	– (100)	– (100)	– (100)
	\mathcal{P}_2	40	†	446.94 (47)	– (100)	– (100)
	\mathcal{P}_2	60	†	884.20 (26)	115.73 (3)	101.91 (3)

– no conv.

* Lower final residual norm than other methods

† Out of Memory

Table: Running time in seconds, and in parenthesis the number of iterations. Stopping tolerance $\text{tol} = 5 \cdot 10^{-6}$. Truncation tolerance $\text{tolrank} = 10^{-12}$. \mathcal{P}_1 : one-term precondition, \mathcal{P}_2 : two-term precondition, expensive.

Numerical example III – [Powell et al. 2017]

Let \mathbf{A}_i and \mathbf{B}_i be the data from [Powell et al. 2017, Example 5.1 and 5.2] coming from the discretization of a 2-dimensional elliptic PDE problem with correlated random inputs,

$$-\nabla \cdot (a(x, \omega) \nabla u) = f \quad \text{in } D, \quad u(x, \omega) = 0, \quad \text{on } \partial D.$$

with $\omega \in \Omega$ a sample space associated with a proper probability space, and $D \subset \mathbb{R}^2$ is the space domain, and

$$a(x, \omega) = \mu(x) + \sigma \sum_{j=1}^{\ell-1} \sqrt{\lambda_j} \phi_j(x) \xi_j(\omega),$$

where μ corresponds to the diffusion coefficient expected value, σ is the standard deviation, while (λ_j, ϕ_j) are the leading eigenpairs of the associated covariance matrix.

The right-hand side has rank 1, while the linear equation counts 10 terms.

Numerical example III – results

Example ($\ell = 10$)	n_A, n_B	maxrank	R-NLCG	MultiRB (spacedim/rank)	TPCG	Ss-CG determ.	Ss-CG rand'zed
[Ex.5.1]	16129, 2002	25 50 100	*5.58 (28) 14.63 (26) 35.98 (25)	6.89 (158/66)	6.55 (24) 13.03 (16) 37.11 (16)	– (100) 4.89 (8) 6.39 (6)	– (100) 3.20 (8) 3.82 (6)
[Ex.5.2]	16129, 1287	60 125 150	– (100) 67.54 (38) 57.31 (24)	12.76 (312/306)	– (100) 53.50 (18) 66.88 (14)	– (100) 19.55 (12) 23.73 (11)	– (100) 11.83 (13) 12.58 (11)

– no conv.

* Final residual norm is *larger* than for other methods

Table: Running time in seconds, and in parenthesis the number of iterations. Stopping tolerance $\text{tol} = 10^{-6}$. Truncation tolerance $\text{tolrank} = 10^{-12}$. Best running times are in bold. For MultiRB the the final approximation space dimension and the final solution rank are reported.

Summary

We presented the Subspace CG, discussing

- matrix coefficients for the CG iterates
- new orthogonality conditions with richer subspaces
- computational enhancements based on randomization and inexactness
- memory costs comparable to TPCG
- promising on numerical examples

Further details on preconditioning, orthogonality and optimality can be found in

D. Palitta, M. Iannacito, and V. Simoncini. [A Subspace–Conjugate Gradient Method for Linear Matrix Equations](#). *SIAM J. Matrix Anal. & Appl.*, 46(4):2197–2225, 2025.



References I

-  Antoulas, A. C. (2005). *Approximation of Large-Scale Dynamical Systems*. Society for Industrial and Applied Mathematics.
-  Benner, P. and T. Breiten (2013). “Low rank methods for a class of generalized Lyapunov equations and related issues”. In: *Numer. Math.* 124, pp. 441–470.
-  Bioli, I., D. Kressner, and L. Robol (2025). “Preconditioned Low-Rank Riemannian Optimization for Symmetric Positive Definite Linear Matrix Equations”. In: *SIAM Journal on Scientific Computing* 47.2, A1091–A1116.
-  Kressner, D. and C. Tobler (2011). “Low-Rank Tensor Krylov Subspace Methods for Parametrized Linear Systems”. In: *SIAM. J. Matrix Anal. & Appl.* 32.4, pp. 1288–1316.
-  Powell, C. E., D. Silvester, and V. Simoncini (2017). “An Efficient Reduced Basis Solver for Stochastic Galerkin Matrix Equations”. In: *SIAM J. Sci. Comput.* 39.1, A141–A163.
-  Shank, S. D., V. Simoncini, and D. B. Szyld (2016). “Efficient low-rank solutions of Generalized Lyapunov equations”. In: *Numerische Mathematik* 134.2, pp. 327–342.

-  [Simoncini, V. \(2016\). “Computational Methods for Linear Matrix Equations”. In: *SIAM Review* 58.3, pp. 377–441.](#)
-  [Simoncini, V. and Y. Hao \(2023\). “Analysis of the Truncated Conjugate Gradient Method for Linear Matrix Equations”. In: *SIAM. J. Matrix Anal. & Appl.* 44.1, pp. 359–381.](#)